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# XIP6110B: ML-KEM 512/768/1024

## Balanced Post-Quantum Key Encapsulation IP Core

Product Brief

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### Introduction

XIP6110B from Xiphera is an Intellectual Property (IP) core for ML-KEM [3] (previously known as CRYSTALS-Kyber [2]) post-quantum Key Encapsulation Mechanism (KEM). It supports key generation, encapsulation, and decapsulation operations for all ML-KEM variants ML-KEM-512, ML-KEM-768, and ML-KEM-1024. XIP6110B is optimized for a good balance between speed and resource requirements.

XIP6110B is a member of xQlave® product family of secure and efficient IP cores for post-quantum cryptography (PQC) algorithms.

### Key Features

- **Small Resource Requirements:** XIP6110B fits into less than 20000 4LUTs and additionally uses only a few multipliers/DSP blocks and internal memory block in a typical Microchip® FPGA implementations.
- **Fast Performance:** XIP6110B is capable of computing a few thousand key generation, encapsulation, or decapsulation operations in a second in a typical Microchip® FPGA implementation.
- **Secure Architecture:** The execution time of XIP6110B is independent of the secret values and, consequently, provides full protection against timing-based side-channel attacks. XIP6110B has been implemented only in digital logic without any software components.
- **Easy Integration:** The simple 64-bit interface of XIP6110B supports easy integration to various systems.
- **Compliance:** XIP6110B is compliant with ML-KEM specifications in NIST FIPS PUB 203 [3].

## Functionality

XIP6110B can be used for key generation, encapsulation, and decapsulation operations of all ML-KEM variants ML-KEM-512, ML-KEM-768, and ML-KEM-1024<sup>1</sup>. ML-KEM was selected as the primary algorithm for post-quantum key encapsulation by the NIST [1] and, hence, it is expected to be very widely used in multiple different protocols in the coming years.

The main optimization objective for XIP6110B has been on achieving a good balance between resource requirements and performance as well as in providing a versatile support for all operations of all ML-KEM variants with a single IP core.

XIP6110B also includes protections against side-channel attacks, the most important of which is that the operation latency does not depend on any secret values. Because ML-KEM uses rejection sampling to obtain certain non-secret values (most importantly, the public key), the actual latency varies slightly between different execution runs. However, because the secret values are obtained in fully constant time by XIP6110B, the variance in the operation latencies does not induce any weaknesses against side-channel attacks.

XIP6110B implements the ML-KEM operations defined in [2], but key generation and encapsulation require random bytes as inputs. Hence, XIP6110B requires an external random number generator (for example, XIP8001B) for generating high-quality random bytes.

## Block Diagram

The internal high-level block diagram of XIP6110B is depicted in Figure 1.

## Interfaces

The external interfaces of XIP6110B are depicted in Figure 2.

This Product Brief describes a high-level overview of the functionality and capabilities of XIP6110B. Please contact sales@xiphera.com for a complete datasheet with a detailed description of the input and output signals, startup procedure of XIP6110B, example simulation waveforms, and the FPGA resource requirements of your targeted FPGA family.

## FPGA Resources and Performance

Table 1 presents the Microchip<sup>®</sup> FPGA resource requirements for representative implementations on different Microchip<sup>®</sup> FPGA architectures. On request, the resource estimates can also be supplied for other Microchip<sup>®</sup> FPGA families.

Microchip <sup>®</sup> Device	Resources	$f_{MAX}$
Microchip <sup>®</sup> PolarFire <sup>®</sup> <sup>¶</sup>	19674 4LUT, 10 LSRAM, 7 Math	121.08 MHz

Table 1: Resource usage and performance of XIP6110B on representative Microchip<sup>®</sup> FPGA families.

<sup>1</sup>“90s” variants of ML-KEM, which are based on AES and SHA-2, that are also defined in [2] are not supported.

<sup>¶</sup>Libero 2022.1.0.10, default compilation settings, industrial speedgrade.

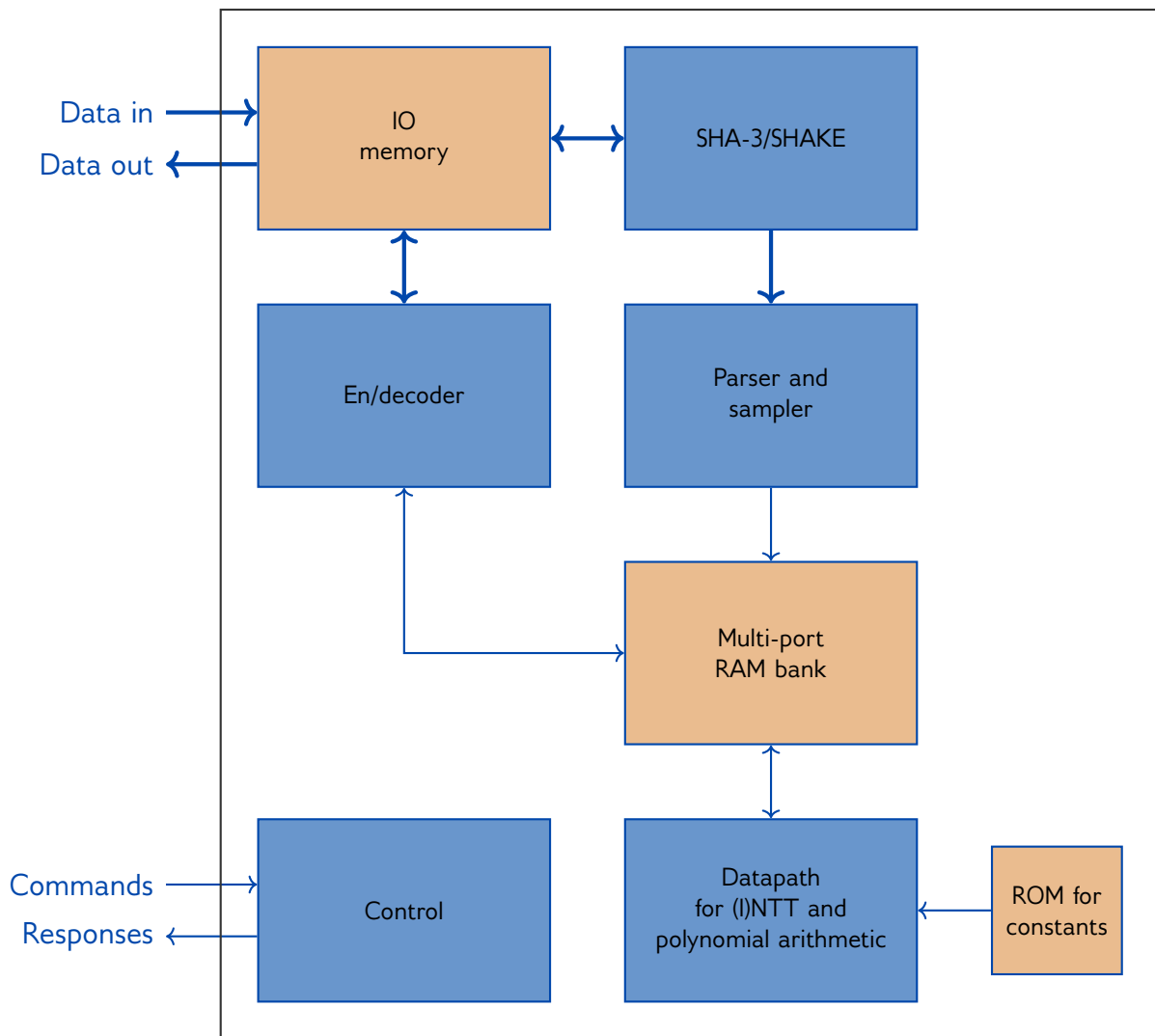


Figure 1: Internal high-level block diagram of XIP6110B

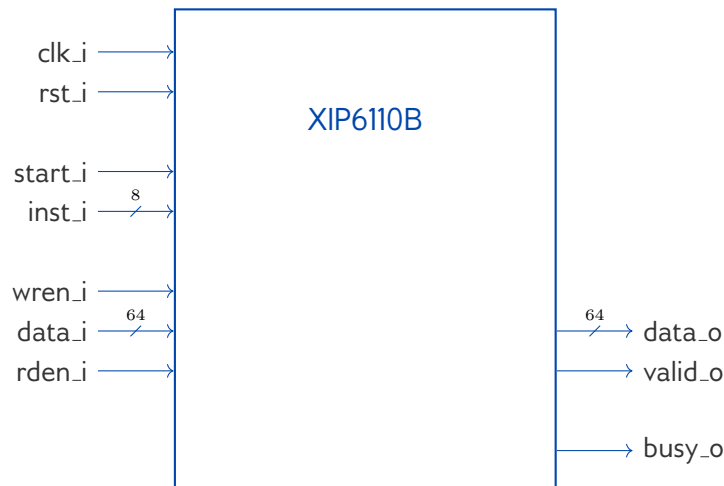


Figure 2: External interfaces of XIP6110B

## Example Use Cases

ML-KEM can be expected to be used for key exchange in various security protocols in the coming years. Therefore, XIP6110B will have several applications in protecting critical systems. There are already drafts about how ML-KEM will be used in security protocols: for example, IPsec IKEv2 [5] and TLS 1.3 [4].

## Ordering and Deliverables

Please contact [sales@xiphera.com](mailto:sales@xiphera.com) for pricing and your preferred delivery method. XIP6110B can be shipped in a number of formats, including netlist, source code, or encrypted source code. Additionally, synthesis scripts, a comprehensive testbench, and a detailed datasheet including an integration guide are included.

## About Xiphera

Xiphera specializes in secure and efficient implementations of standardized cryptographic algorithms on Field Programmable Gate Arrays (FPGAs) and Application Specific Integrated Circuits (ASICs). Our fully in-house designed product portfolio includes individual cryptographic Intellectual Property (IP) cores, as well as comprehensive security solutions built from a combination of individual IP cores.

Xiphera is a Finnish company operating under the laws of the Republic of Finland, and is fully owned by Finnish citizens and institutional investors.

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## References

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